A Forward-Looking Nash Game and Its Application to Achieving Pareto-Efficient Optimization

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ABSTRACT

Recognizing the fact that a player’s cognition plays a defining role in the resulting equilibrium of a game of competition, this paper provides the foundation for a Nash game with forward-looking players by presenting a formal definition of the Nash game with consideration of the players’ belief. We use a simple two-firm model to demonstrate its fundamental difference from the standard Nash and Stackelberg games. Then we show that the players’ belief functions can be regarded as the optimization parameters for directing the game towards a much more desirable equilibrium.

Keywords: Belief; Cognition; Iterative Algorithm; Nash Equilibrium; Pareto-Optimality; Stackelberg

1. Introduction

Game theory has been very well recognized as a branch of applied mathematical tools, best for analyzing the phenomenon of selfish competition, which arises in numerous real-life applications ranging from economics to social sciences, and even to engineering problems. Many optimization problems can be viewed as a competition problem.

If there are shared resources, then inherently there will be competition in the allocation of such resources. An individual who participates in the resource allocation always wishes to maximize its payoff. Nevertheless, the player’s return not only depends on its own strategy, but is also dependent on competitors’ responses. As a result, ideally, a player should choose or optimize its strategy based on not only its immediate return but also the possible outcomes of how others might respond to its strategy. Apparently, a player’s cognition (i.e., its belief on how the environment as a whole would react to any of its actions) will be pivotal in the optimization of its strategy for maximizing its payoff and in defining the equilibrium of the game.

In the literature, there are a number of game-theoretic models and they differ in their assumptions on the players’ cognition. Due to this fundamental difference, these models represent different competition environments, thereby resulting in very different outcomes, which at the steady state are referred to as equilibria. When the competition process reaches an equilibrium, by definition, no player would have incentive to change their strategies further to deviate from the equilibrium which is therefore usually considered as a satisfactory outcome to all players.

The most popular model has to be the Nash game [1] in which a player has the belief that other players’ strategies are fixed and will not change regardless of what it does. A Nash game is based on such assumption on the players’ cognition. It is noted that the game is still very well defined, although the player’s belief is totally inaccurate.

On the other hand, in a Stackelberg game [2], there is a super player, commonly known as leader, who knows all the information about its competitors, known as followers. In this case, the leader is considered to have perfect cognition and therefore can obtain the most rewarding strategy.

In most practical competition situations, the efficiency of a Nash game resulting in the most celebrated Nash equilibrium is often poor (which will be revealed by a two-firm example later in this paper). The main reason is that the assumption for the Nash game is too much simplified, and it fails to account for the interaction between
the players. In contrast, the assumption for the Stackelberg model is too stringent and it is often too difficult to be qualified to be a leader in most practical scenarios. Nevertheless, the Stackelberg equilibrium is highly beneficial to the leader. The limitation is that if there were two or more leaders, all of which possess perfect cognition and wish to be the biggest winner, this would lead to a tragedy [3].

In this paper, we recognize the importance of players’ cognition in a game, and focus on how the assumption of the players’ cognition affects the equilibrium. Based on the analysis of Nash and Stackelberg equilibria, we provide a new definition for the Nash game with consideration of the players’ belief. As a useful byproduct, the new definition facilitates the interpretation of the players’ cognition as the optimization parameters, which, if optimized properly, can direct the game towards a Pareto-efficient equilibrium.

2. Preliminaries

2.1. The Two-Firm Model

For illustrative purpose, in this paper, we use a simple two-firm model in economics as an example. In this simple model, there are two firms, Firm A and Firm B. They manufacture an identical product and its unit cost is the same for both firms. In addition, the profit per unit diminishes as the number of products available in the market increases. In particular, let \( x_a \), \( x_b \) denote the number of products manufactured by Firm A, and \( x_b \) be that manufactured by Firm B. The unit cost is denoted by \( c \) and the price per product which is set the same by both Firm A and Firm B is assumed to be \( a - x_a - x_b \) for some constant \( a > 0 \). Therefore, the profit functions for Firm A and Firm B are, respectively, given by

\[
\begin{align*}
    f_a(x_a, x_b) &= (a - x_a - x_b - c)x_a, \\
    f_b(x_a, x_b) &= (a - x_a - x_b - c)x_b.
\end{align*}
\]

(1)

In various game-theoretic models, the objective would be for Firm A and Firm B to iteratively optimize their respective strategies, \( x_a \) and \( x_b \), for maximizing their profits (1) in a competitive fashion. Before we examine different equilibria of this two-firm game, we find it useful to first present the general notations of a game.

For a game with \( K \) players, we denote the strategy profile for player \( k \) as \( S_k \) and use \( S = S_1 \times S_2 \times \cdots \times S_K \) to represent the strategy profile for all the players. Moreover, we use \( x = (x_1, x_2) \in S \) to denote a specific choice of strategy from all the players, where \( x_k \in S_k \) is the strategy adopted by player \( k \), and

\[
x_k = (x_{k,1}, x_{k,2}, \ldots, x_{k,K})
\]

(3)
denotes the adopted strategy by all the players except player \( k \). Likewise, the reward function for player \( k \) is denoted by \( f_k(x) \). If the game converges to an equilibrium, then we will use the superscript \( (\cdot) \) to highlight that the corresponding parameters are at the equilibrium. For instance, we have \( x^* \) and \( f^*_k \) at the equilibrium.

2.2. Nash Competition

Given the competition of \( K \) players, if at some point, no one can gain any further by deviating from its present strategy, then the strategy of all the players is said to have reached to an equilibrium [4]. Mathematically, we have

\[
\forall k \text{ and } \forall x_k \in S_k : f_k(x^*) \geq f_k(x_k, x^*).
\]

(4)

Now, proceed to derive the Nash equilibrium for the two-firm example. According to (4), Firm A solves

\[
\max_{x_a \geq 0} f_a(x_a, x_b) = \max_{x_a \geq 0} (a - x_a - x_b^* - c)x_a.
\]

(5)

As such, it can be easily shown that

\[
\frac{df_a}{dx_a} = a - x_b^* - c - 2x_a \Rightarrow x_a^* = \frac{a - x_b^* - c}{2}.
\]

(6)

Similarly, \( x_b^* \) can be derived, resulting the set of simultaneous equations

\[
\begin{align*}
    x_a^* &= \frac{a - x_b^* - c}{2}, \\
    x_b^* &= \frac{a - x_a^* - c}{2}.
\end{align*}
\]

(7)

As a result, the game governed by (7) will reach the Nash equilibrium

\[
x^* = \left( \frac{a-c}{3}, \frac{a-c}{3} \right).
\]

(8)

which leads to the profits

\[
f^* = \left( f_a^*, f_b^* \right) = \left( \frac{(a-c)^2}{9}, \frac{(a-c)^2}{9} \right).
\]

(9)

It is worth pointing out here that in Nash equilibrium when taking \( \frac{df_a}{dx_a} \) in the optimization for player A, player A believes that \( x_b \) at any given time instance is optimal and fixed, or \( \frac{dx_b}{dx_a} = 0 \). This is obviously inaccurate. At the Nash equilibrium in particular, from (1.7), we actually have \( \frac{dx_b}{dx_a} = -0.5(\neq 0) \). This shows
that a Nash player’s cognition and the reality have considerable difference and the profit function \( f_a(x_a, x'_b) \) is in fact only the profit in the ideal situation \( x_a = x'_b \) but does not represent the actual profit function due to interaction from player B.

### 2.3. Stackelberg Competition

A player can benefit more from the rest of the players in a game if this player has perfect cognition about how others would react to its strategy. This is studied formally by the Stackelberg equilibrium. In the general setting with \( K \) players, if player \( \kappa \) is the leader, it should know perfectly the response function \( x_a(x_b) \) and therefore is able to obtain the most effective strategy by the Stackelberg equilibrium. In the general setting \( K \) players, if player \( \kappa \) is the leader, it should know perfectly the response function \( x_a(x_b) \) and therefore is able to obtain the most effective strategy such that

\[
\forall x_{\kappa} \in S_{\kappa}: f_{\kappa}(x^*) \geq f_{\kappa}(x_\kappa, x_{-\kappa}(x_\kappa)). \tag{10}
\]

For other players \( k \neq \kappa \), as they do not have any information about any of the other players’ strategies, they can only assume that other players’ strategies are fixed, and will act like a Nash player, see (4). As a result, we have

\[
\forall k \neq \kappa \text{ and } \forall x_k \in S_k: f_k(x^*) \geq f_k(x_k, x^*_a). \tag{11}
\]

In a Stackelberg game, there is a very strict order of how the players play the game [1,5]. In particular, leader \( \kappa \) needs to first give out its strategy and lets other players compete to reach a Nash equilibrium against this strategy before it revises its strategy for another round of competition among the rest of the players. Achieving the Stackelberg equilibrium will thus require a two-level game.

The merit of Stackelberg equilibrium is that leader \( \kappa \) has an absolute advantage over other players but the drawback is that knowing the function \( x_a(x_b) \) would be too difficult to achieve in practice, if not impossible. A standard approach would require the leader to try exhaustively all possible \( x_{\kappa} \in S_{\kappa} \) to identify the best strategy.

Recalling from the two-firm example, if we let Firm A be the leader and Firm B be the follower, then according to (11), player B’s cognition is that player A’s strategy is optimal and fixed and player B therefore aims to solve

\[
\max_{x_b \geq 0} f_b(x_b, x^*_a) = \max_{x_b \geq 0} \left( a - x^*_a - x_b - c \right) x_b, \tag{12}
\]

which by setting \( \frac{df_b}{dx_b} = 0 \) gives

\[
x_b = \frac{a - x^*_a - c}{2}. \tag{13}
\]

which is the exact response function of player B with respect to any action \( x_a \).

According to (10) and using (13), player A finds its best strategy by

\[
\max_{x_a \geq 0} f_a(x_a, x^*_b) = \max_{x_a \geq 0} \left( a - x_a - \frac{a - x_a - c}{2} \right) x_a. \tag{14}
\]

As a result, the best strategy for player A can be analytically obtained as

\[
x^*_a = \frac{a - c}{2} \tag{15}
\]

and the corresponding strategy for player B is

\[
x^*_b = \frac{a - c}{4}. \tag{16}
\]

Hence, at the Stackelberg equilibrium, we have

\[
x^* = \left( \frac{a - c}{2}, \frac{a - c}{4} \right), \tag{17}
\]

which leads to the profits

\[
f^* = \left( \frac{(a-c)^2}{8}, \frac{(a-c)^2}{16} \right). \tag{18}
\]

Note that in order to achieve the Nash equilibrium earlier, there is no specific order of obtaining \( x^*_a \) and \( x^*_b \) in solving the simultaneous Equations (7) and the equilibrium can be achieved by free competition. This is however not true for the Stackelberg equilibrium where the leader’s strategy, \( x^*_a \), must be obtained first and the follower(s) respond. As the number of players increases, the complexity of the Stackelberg game will increase considerably and the simplicity of the Nash game will prevail. In addition, a game with all leaders degenerates to a Nash game.

In the two-firm example, because Firm A knows precisely the strategy adopted by Firm B (13), it is able to achieve a higher profit than what is achieved by the Nash equilibrium. However, the key questions are:

- If (13) is not known by Firm A, or Firm A only knows partial information about (13), e.g., \( dx_a / dx_b \) or only \( dx^*_b / dx^*_a \), would it still help Firm A to obtain a better or even Stackelberg strategy? If the answer is yes, this would mean that the level of cognition for a leader could be significantly reduced.
- Also, what if two players have partial information about each other’s strategies? What will happen?

This has motivated us to investigate the Nash game with simultaneous forward-looking players, in which each player optimizes its strategy based on its own belief.

### 3. Forward-Looking Competition

After the discussion of Nash and Stackelberg equilibria above, it is clear that players’ cognition is key to defining the resulting equilibrium of a game. At the same time, we
understand the beauty of the Nash game where players can compete freely (without following a specific order) to reach the equilibrium. Based on these two points, we present the Nash game with forward-looking players.

To facilitate our analysis, we find it useful to have the following definitions.

- **Environmental function**—In a competition process, the reward for player $k$ depends not only its own strategy $x_k$ but also others’ strategies $x_{-k}$, at any given time instant $t$. We use the environmental function $r_i(x', x_{-k})$ to quantify the influence of other players’ strategies (at time instant $t$) onto player $i$’s reward. Obviously, other players’ strategies can always be treated as some form of response to a given player’s strategy at time instant $t$.

- **Belief function**—A player’s understanding on its environmental function reflects its cognition about the competition in the game. We assume that player $k$ possesses the knowledge of a belief function, which is denoted as $r_k^b(x_k', x_{-k})$, where $x_{-k}$ denotes the strategies from all the players at time instant $t$ except player $k$, and clearly $r_k^b(x_k', x_{-k}) = r_k(x_{-k})$. The latter relationship further suggests to formulate the belief response using some form of Taylor series expansion. For example, we may write

$$r_k^b(x_k, x_{-k}) = r_k(x_{-k}) + \phi_k^b(x_k - x_k'),$$  

(18)

where $\phi_k^b \triangleq \frac{d r_k^b(x_k, x_{-k})}{dx_k}$ is regarded as the interference derivative (to be discussed in Section 3.3). The belief function $r_k^b(\cdot)$ is player $k$’s cognition on what the environmental function $r_k(\cdot)$ would be, given a strategy $x_k$ and the present state $x_{-k}$. If $r_k^b(\cdot) = r_k(\cdot)$, then player $k$’s cognition will be perfect. Otherwise, it is only a prediction.

- **Predicted reward**—The predicted reward function, denoted by $f_i(x, r_i^b(x_k, x_{-k}))$, indicates the amount of reward player $k$ believes to achieve by the strategy $x_k$ and other players’ strategies at time instant $t$, $x_{-k}$, based on the belief function $r_i^b(\cdot)$.

### 3.1. A Motivating Example

In Section 2.3, if player A is a Stackelberg leader, then it has the environmental function (or response) $r_A = x_b$ and knows perfectly the strategy of player B (13). Therefore, it has the perfect belief function

$$r_A^b(x_b) = r_A = x_b = 0.5(a - x_a - c).$$  

(19)

If player A’s cognition is reduced; for instance, knows at any given time instant $\frac{dx_b}{dx_b} = \frac{dx_b}{dx_A} = -0.5$ in (13) and can observe the environmental function at present time $t$, i.e., $r_A^b(x'_b, x'_b) = r_A(0.5a - x'_a - c)$, player A can formulate its belief function, using a first-order Taylor series, as

$$r_A^b(x_b, x'_b) = x'_b - 0.5(x'_b - x_b).$$  

(20)

and its predicted reward function can be expressed as

$$f_A(x_b, r_A^b(x_b, x'_b)) = \left(a - x_b - r_A^b - c\right)x_b$$

$$= \left(a - x'_b - 0.5x'_b - c - 0.5x_b\right)x_b.$$  

(21)

From the side of player B, if it has as low cognition as a Nash player, then it will have the belief function $r_B^b = x'_b$. Hence, player B’s reward function will be

$$f_B(x_b, r_B^b(x_b, x'_b)) = \left(a - r_B^b - x_b - c\right)x_b$$

$$= \left(a - x'_b - c - x_b\right)x_b.$$  

(22)

As a result, player A’s strategy can be optimized by

$$\begin{align*}
\max_{x_a \geq 0} f_A(x_a, r_A^b(x_a, x'_b)) &= \max_{x_a \geq 0} \left(a - x'_b - 0.5x'_b - c - 0.5x_b\right)x_a.
\end{align*}$$  

(23)

By setting $\frac{df_A^b}{dx_a} = 0$, this suggests an updating process

$$x'_{a+1} = a - x'_b - 0.5x'_b - c.$$  

(24)

On the other hand, for player B, it aims to solve

$$\begin{align*}
\max_{x_b \geq 0} f_B(x_b, r_B^b(x_b, x'_b)) &= \max_{x_b \geq 0} \left(a - x'_b - c - x_b\right)x_b.
\end{align*}$$  

(25)

By setting $\frac{df_B^b}{dx_b} = 0$, player B updates its strategy by

$$x'_{b+1} = a - x'_b - c.$$  

(26)

Applying the two updating rules (24) and (26) repeatedly, as $t \to \infty$, we get

$$\begin{align*}
x' &= \left(a - c - a - c\right) \frac{2}{4},
\end{align*}$$  

(27)

which ends up the same strategy in the Stackelberg equilibrium (16), but via a completely different process.

The striking result here is that the Stackelberg equilibrium can now be achieved by a free competition between the players without following a specific order of how the game should be played and that the so-called leader, i.e., player A here, does not need perfect cognition of the environmental function (13), but a good belief (20).

### 3.2. Definition of the New Equilibrium

Motivated by the potential of being forward-looking, we here present the Nash equilibrium with forward-looking
players (i.e., with some cognition in the form of belief functions). Mathematically, it is written as

$$\forall k \text{ and } x_k \in S_k :$$

$$f_k \left( x_k^*, r_k^\alpha \left( x_k, x_k^* \right) \right) \geq f_k \left( x_k, r_k^\alpha \left( x_k, x_k^* \right) \right),$$

(28)

where \( r_k^\alpha (\cdot) \) is the belief function reflecting player \( k \) ’s cognition ability and \( f_k (\cdot) \) is the predicted reward function for player \( k \). Note that (28) can be rewritten as

$$\forall k \text{ and } x_k \in S_k : f_k \left( x_k^*, r_k^\alpha \left( x_k, x_k^* \right) \right) \geq f_k \left( x_k, r_k^\alpha \left( x_k, x_k^* \right) \right),$$

(29)

because according to (18), \( r_k^\alpha \left( x_k, x_k^* \right) = r_k \left( x_k^* \right) \) and as such, at the equilibrium, we have

$$f_k \left( x_k^*, r_k^\alpha \left( x_k, x_k^* \right) \right) = f_k \left( x_k^*, r_k \left( x_k^* \right) \right) = f_k \left( x^* \right).$$

(30)

In this model, the belief function \( r_k^\alpha (\cdot) \) can be chosen arbitrarily and it only serves to indicate player \( k \) ’s understanding about the competition environment. In fact, (29) embraces the conventional Nash equilibrium (4) in which players have the belief function \( r_k^\alpha \left( x_k, x_k^* \right) = r_k \left( x_k^* \right) \) which effectively treats the environment player \( k \) observes at any present time instant \( t \) as fixed and constant and ignores the subsequent changes in other players’ strategies provoked by player \( k \)’s new strategy.

We refer to the equilibrium of a Nash game with belief functions as a belief-directed Nash equilibrium (BNE).

### 3.3. From Interference Derivative to Pareto-Optimality

In the proposed model of the Nash game with forward-looking players, every player is free to choose its own belief function. Altogether, the combination of the players’ belief functions defines the resulting equilibrium of the players’ competition and has numerous possibilities. To examine this, we consider in the two-firm example that

$$r_a^\alpha (x_a, x_a^*) = x_a^* + \lambda_a (x_a - x_a^*),$$

$$r_b^\alpha (x_b, x_b^*) = x_b^* + \lambda_b (x_b - x_b^*),$$

(31)

where \( \lambda = (\lambda_a, \lambda_b) \) are regarded as the interference derivatives which can be interpreted as the rate of change of the environmental function with respect to one’s strategy. Also, (31) can be viewed as a first-order Taylor series approximation for the response. With (31), we can express the predicted reward functions for the players as

$$f_a \left( x_a, r_a^\alpha (x_a, x_a^*) \right) = \left( a - x_a^* - c - \lambda_a (x_a - x_a^*) - x_a \right) x_a,$$

$$f_b \left( x_b, r_b^\alpha (x_b, x_b^*) \right) = \left( a - x_b^* - c - \lambda_b (x_b - x_b^*) - x_b \right) x_b.$$

(32)

For convenience, we refer to this two-firm example as a BNE by BNE(\( \lambda \)). In this game, player A aims to

$$\max_{x_a \geq 0} f_a \left( x_a, r_a^\alpha \left( x_a, x_a^* \right) \right) = \max_{x_a \geq 0} \left( a - x_a^* - c - \lambda_a (x_a - x_a^*) - x_a \right) x_a,$$

(33)

which can be solved by setting \( \frac{dx_a^*}{dx_a} = 0 \). This then gives

$$x_a^{*1} = \frac{a - x_a^* + \lambda_a x_a - c}{2(\lambda_a + 1)}.$$

(34)

Similarly, we can easily derive the corresponding result for optimizing \( x_b \). As \( t \to \infty \) (after a sufficient number of iterations), at the equilibrium, we have

$$\left\{ \begin{array}{ll}
x_a^* = \frac{(a-c)(\lambda_a + 1)}{(\lambda_a + 2)(\lambda_a + 2) - 1}, \\
x_b^* = \frac{(a-c)(\lambda_b + 1)}{(\lambda_b + 2)(\lambda_b + 2) - 1}.
\end{array} \right.$$

(35)

The above strategies will lead to the profits (or rewards)

$$f_a^* = \frac{(a-c)^2(\lambda_a + 1)(\lambda_a + 1)^2}{[\lambda_a + 2)(\lambda_a + 2) - 1]^2},$$

$$f_b^* = \frac{(a-c)^2(\lambda_b + 1)(\lambda_b + 1)}{[\lambda_b + 2)(\lambda_b + 2) - 1]^2}.$$

(36)

Very Interestingly, we can observe that the equilibrium varies according to the belief parameters of the players, \( \lambda \), which offers an opportunity to optimize the equilibrium.

In order to illustrate this, we let \( \lambda_a = \lambda_b = \lambda \) and consider that \( \lambda \) is an optimization parameter of the game. Then \( \lambda \) can be optimized by

$$\max_{\lambda \geq -3} \frac{(a-c)^2(\lambda + 1)}{(\lambda + 2)^2},$$

(37)

The above optimization maximizes both the sum-profit and the individual profits of the players, and therefore is Pareto-optimal. Note that \( \lambda = -3 \) is not permitted as this would lead to unbounded solutions to \( x_a^* \) and \( x_b^* \).

It can be easily shown that the optimal value of \( \lambda \) is 1, which gives

$$x_a^* \bigg|_{\lambda=1} = \left( \frac{a-c}{4}, \frac{a-c}{4} \right),$$

(38)

and

$$f_a^* \bigg|_{\lambda=1} = \left( \frac{(a-c)^2}{8}, \frac{(a-c)^2}{8} \right).$$

(39)
We summarize our results and include some interesting scenarios in Table 1. Note that in this two-firm example, the Stackelberg equilibrium is not Pareto-optimal. It can be seen by comparing it with BNE(1) that if the leader is being less aggressive, its profit is not reduced but the profit for the follower can be further increased.

Another observation is that achieving Pareto-optimum does not require that the players’ cognition be accurate. In particular, for the case of Pareto-optimum, i.e., \( \lambda = 1 \), we see from (34) that at the equilibrium, we have

\[
\frac{dr_a^*}{dx_a} = \frac{dx_a^*}{dx_a} = -3. \tag{40}
\]

However, intriguingly, the player’s cognition is very different and for player A, we have

\[
\frac{dr_a^b}{dx_a} = \lambda = 1 \neq -3. \tag{41}
\]

This reveals that the belief function a player has is not required to reflect the true reality but can still achieve the best outcome from the competition of the players. This has changed our understanding on how player’s cognition influences the equilibrium of a game. A beautiful world may be an outcome from people’s mistakes.

On the contrary, if \( \lambda = -1 \), then it can be shown that

\[
\frac{dr_a^*}{dx_a} = \frac{dx_a^*}{dx_a} = -1 \text{ and } \frac{dr_b^*}{dx_b} = \frac{dx_b^*}{dx_b} = -1, \tag{42}
\]

and the belief functions are perfect in terms of representing the reality. Unfortunately, the profits for both players at the equilibrium will be 0, showing that if both players are perfectly smart, the outcome could be a tragedy.

In this table, we also included the case with \( \lambda = 3 \) which interestingly results in the same profits as the Nash equilibrium but with half production. To show how the belief (or the interference derivative \( \lambda \)) affects the strategy \( x^* = x_a^* = x_b^* \) and the reward \( f^* = f_a^* = f_b^* \), we provide the results in Figure 1 assuming that \( a - c = 10 \).

### 3.4. Existence

Since the birth of game theory, the quest for the existence of an equilibrium has been the focal point in this area. In [6], Nash proved the existence of the widely known Nash equilibrium we know today. Later, numerous researchers presented further theorems and proofs about the existence of the game-theoretic equilibria [7,8]. Based on what is known in the literature, we provide the sufficient condition for the existence of the general BNE game.

**Corollary 1** Consider a K-player game, with the strategy spaces \( S_k \) for \( k = 1, 2, \ldots, K \), which are not empty, compact and convex subsets of an Euclidean space. If the predicted reward function \( f_k(x, r_k^b(x_k, x_{-k})) \forall k \) is continuous and quasi-concave in \( S_k \) for any fixed \( x_{-k} \), then there exists an equilibrium.

![Figure 1. The strategy and reward at various equilibria against the belief \( \lambda \) assuming \( a - c = 10 \). The solid line refers to the results for the strategy while the dash line shows the rewards.](Image)

\[\text{Table 1. Strategies and profits for various equilibria.}\]

<table>
<thead>
<tr>
<th>Model</th>
<th>Strategy ((x^*))</th>
<th>Profit/Reward ((f^*))</th>
</tr>
</thead>
<tbody>
<tr>
<td>Nash Equilibrium</td>
<td></td>
<td></td>
</tr>
<tr>
<td>BNE((\lambda)) with (\lambda = (0, 0))</td>
<td>(2 \left(\frac{a - c}{3}, \frac{a - c}{3}\right))</td>
<td>(3 \left(\frac{(a - c)^2}{9}, \frac{(a - c)^2}{9}\right))</td>
</tr>
<tr>
<td>BNE((\lambda)) with (\lambda = (3, 3))†</td>
<td>((\frac{a - c}{6}, \frac{a - c}{6}))</td>
<td>((\frac{(a - c)^2}{16}, \frac{(a - c)^2}{16}))</td>
</tr>
<tr>
<td>Stackelberg Equilibrium</td>
<td></td>
<td></td>
</tr>
<tr>
<td>BNE((\lambda)) with (\lambda = (-0.5, 0))†</td>
<td>((\frac{a - c}{4}, \frac{a - c}{4}))</td>
<td>((\frac{(a - c)^2}{8}, \frac{(a - c)^2}{8}))</td>
</tr>
<tr>
<td>BNE((\lambda)) with (\lambda = (1, 1))†</td>
<td>((\frac{a - c}{4}, \frac{a - c}{4}))</td>
<td>((\frac{(a - c)^2}{8}, \frac{(a - c)^2}{8}))</td>
</tr>
<tr>
<td>BNE((\lambda)) with (\lambda = (-1, -1))</td>
<td>((\frac{a - c}{2}, \frac{a - c}{2}))</td>
<td>((0, 0))</td>
</tr>
</tbody>
</table>

†It has the same profits as Nash equilibrium but half production. *It achieves the Stackelberg equilibrium but is not Pareto-optimal. †This is the Pareto-optimal equilibrium.
Proof The BNE game (with forward-looking players) presented in Section 3.2 follows the same spirit as a typical Nash game where any player acts towards an equilibrium assuming that other players’ strategies are fixed. Their only difference lies in their understanding about the payoff functions due to different level of cognition to the environment. Consequently, the same result regarding existence of Nash equilibrium [7] can be directly applied to BNE by simply replacing the pay off function by the predicted reward function, which completes the proof.

4. Conclusion

This paper showed the importance of players’ cognition to the equilibrium of the game and presented the Nash equilibrium of forward-looking players. Using a two-firm example, we demonstrated how a Nash game with belief (referred to as BNE) can be made to achieve the Nash and Stackelberg equilibria. On the other hand, this paper has also illustrated the potential of using the belief function as an optimization parameter which makes possible the game converging to a Pareto-optimal equilibrium. However, this is a new regime in game theory and future work is required to better understanding of belief-directed games.

REFERENCES